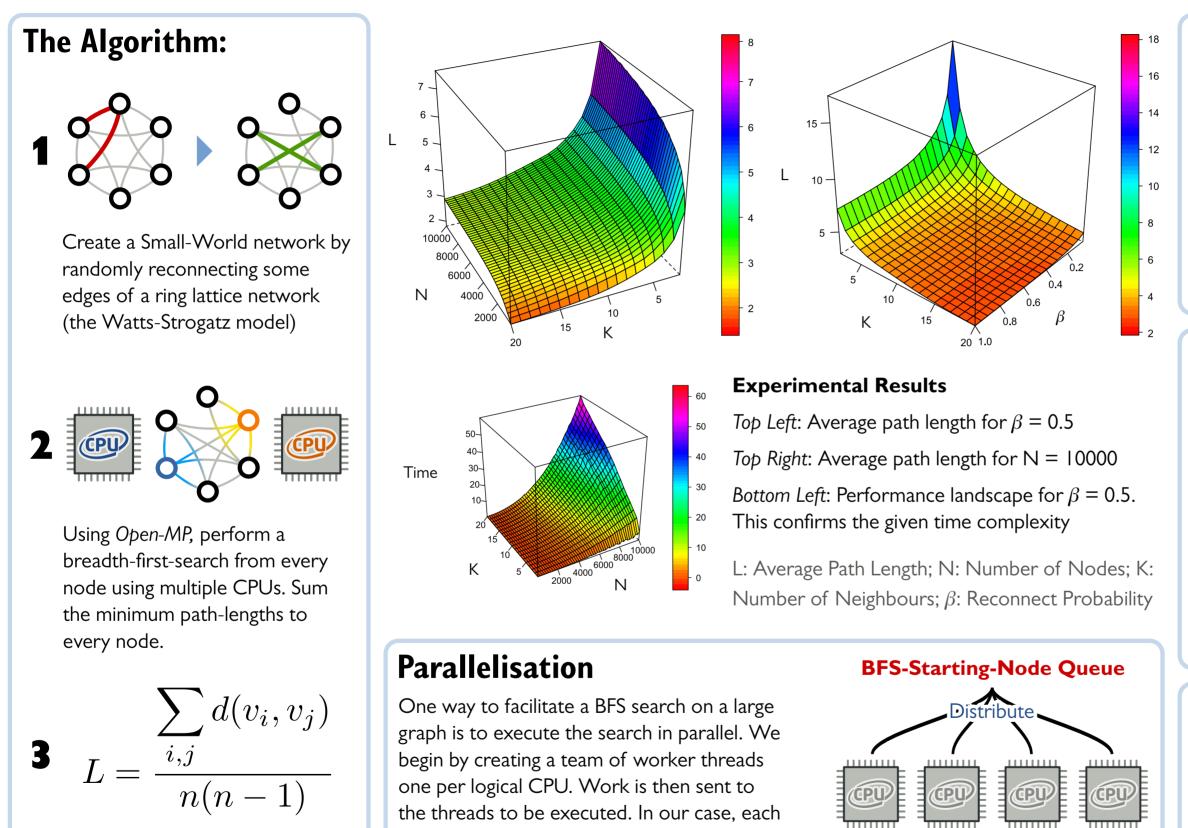
AVERAGE PATH LENGTH IN SMALL-WORLD NETWORKS

Christopher David Williams • Jonathan Heathcote • Karl Sutt • Matt Leach • Tom Nixon



thread executes a BFS origininating from a

different node. All threads share the same

graph structure in memory but have an

individual search state.

Divide the summed shortest-path lengths by the number of paths to find the average path length.

Vs. Floyd's Algorithm

Performance determined by the number of vertices in the graph running in $O(V^3)$. This experiment uses a small-world network - this has, by definition, a large number of vertices with a relatively low number of edges between them and so this is not helpful for this algorithm.

Vs. Dijkstra's Algorithm

Dijkstra's algorithm outperforms Floyd's due to the significant term being $O(V \log, V)$ not $O(V^3)$. An everysource dense graph problem is solved in $O(\frac{1}{2}V^2(V + 2\log V - 1))$ - marginally faster than Floyd's. In practice, BFS suits the given problem better, as edge weights are constant and complexity being O(EV). Since E = KN, V = N, we have $O(KN^2)$.

References

Reduce

Average Path Length

Duncan J. Watts & Steven H. Strogatz; "Collective dynamics of 'small-world' networks" Nature Volume 393, 4 June 1998. pp 440-442.

OpenMP Architecture Review Board; "**OpenMP Application Program Interface**" Version 3.1, July2011