## AVERAGE PATH LENGTH IN SMALLWMORLD NETWORKS

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## The Algorithm:



Create a Small-World network by randomly reconnecting some edges of a ring lattice network (the Watts-Strogatz model)
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Using Open-MP, perform a breadth-first-search from every node using multiple CPUs. Sum the minimum path-lengths to every node.

$$
3 L=\frac{\sum_{i, j} d\left(v_{i}, v_{j}\right)}{n(n-1)}
$$

Divide the summed shortest-path lengths by the number of paths to find the average path length.


## Parallelisation

One way to facilitate a BFS search on a large graph is to execute the search in parallel. We begin by creating a team of worker threads one per logical CPU. Work is then sent to the threads to be executed. In our case, each thread executes a BFS origininating from a different node. All threads share the same graph structure in memory but have an individual search state.


## Vs. Floyd's Algorithm

Performance determined by the number of vertices in the graph running in $O\left(V^{3}\right)$. This experiment uses a small-world network - this has, by definition, a large number of vertices with a relatively low number of edges between them and so this is not helpful for this algorithm.

## Vs. Dijkstra's Algorithm

Dijkstra's algorithm outperforms Floyd's due to the significant term being $O(V \log , V)$ not $O\left(V^{3}\right)$. An everysource dense graph problem is solved in $O\left(\frac{1}{2} V^{2}(V+2 \log V-1)\right)$ - marginally faster than Floyd's. In practice, BFS suits the given problem better, as edge weights are constant and complexity being $O(E V)$. Since $E=K N, V=N$, we have $O\left(K N^{2}\right)$.

## References

Duncan J. Watts \& Steven H. Strogatz;
"Collective dynamics of 'small-world'
networks" Nature Volume 393, 4 June 1998. pp 440-442.

OpenMP Architecture Review Board; "OpenMP Application Program Interface" Version 3.I, July201I

